

Widely applied math seminar Harvard University Nov. 1, 2010

General packing problem

Let $\phi_{max}(K)$ be the highest achievable density for packings of convex d-dimensional body K.

2D periodic examples:



 $\phi_{max}(K)$ for d > 2 known only for spheres, space-filling solids.

From Hilbert's 18th problem:

"How can one arrange most densely in space an infinite number of equal solids of a given form, e.g., *spheres* with given radii or *regular tetrahedra* with given edges (or in prescribed position), that is, how can one so fit them together that the ratio of the filled to the unfilled space may be as large as possible?"



David Hilbert (1862-1943)



Rogers bound



 $\phi(B) \leq 0.7796$

Argument: the sphere cannot fill its Voronoi region, a polyhedron

Rogers bound



 $\varphi(B) \leq 0.7796$

Argument: the sphere cannot fill its Voronoi region, a polyhedron

The tetrahedron can easily fill its Voronoi region

Regular tetrahedra do not fill space



Missing angle: 7.4°

Therefore, $\phi(T) < 1$

But can we find $\phi^{\cup} < 1$ such that $\phi(T) \le \phi^{\cup}$?

Tetrahedron packing upper bound

Optimization challenge:

1. Prove $\phi \leq 1 - \epsilon$, where $\epsilon > 0$

2. Maximize ε

Tetrahedron packing upper bound

Optimization challenge:

1. Prove $\phi \leq 1 - \epsilon$, where $\epsilon > 0$

2. Maximize ε

2'. Minimize length of proof

Solution: $\epsilon = 2.6... \times 10^{-25}$ (15 pages)

Gravel, Elser, & Kallus, Discrete and Computational Geometry (2010)

Bound from angle mismatch





If we can put a lower bound the amount of uncovered space in a unit ball with five non-overlapping wedges, we can get a non-trivial upper bound on the density of a tetrahedron packing. If all wedge edges pass through the center, we can easily calculate the uncovered volume. Unfortunately, this isn't given.

Still, if all wedge edges pass within a given distance of the center, we can still easily calculate a bound on the uncovered volume.

> And if one or more edge wedges fall outside the yellow sphere, we are left with a simpler configuration inside the yellow sphere, and we can try to put a bound on the uncovered volume inside it.

¹¹ By applying this argument iteratively: $\phi(T) \le 1 - (2.6...) \times 10^{-25}$



Lower bounds (densest known packings)

1. Conway & Torquato (2006)



Icosahedral packing: $\varphi = 0.7166$ N=20 "Welsh" packing: $\varphi = 0.7175$ N=34

Conjecture: $\phi_{max}(T) < \phi_{max}(B)$

PNAS (2006)

Densest known packings

2. Chen (2008)



"wagon wheels" packing: $\phi = 0.7786$ N=18

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\phi_{\max}(T) > \phi_{\max}(B)
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Discrete & Compu. Geom. (2008)

Densest known packings

3. Torquato & Jiao (2009)

Challenge for numerical search: highly frustrated optimization problem

Search often got stuck at local optima



Nature (2009)





 $\phi = 0.8226$ N = 314

Phys. Rev. E (2009)

Densest known packings

4. Kallus et al. (2009)

Method: "divide and concur"



 $\phi = 0.8547$ N = 4 (!)

All tetrahedra equivalent (tetrahedron-transitive packing)

Discrete & Compu. Geom. (2010)

5. Chen et al. (2010)

Slight analytical improvement to the above structure: $\phi = 0.8563$ (New, denser, packing is no longer tetrahedron-transitive)

Discrete & Compu. Geom. (2010)

Computational approach to packing problems:

Optimization: given a collection of figures, arrange them without overlaps as densely as possible.

Feasibility: find an arrangement of density $> \phi$

Possible approaches:

- Complete algorithm
- Specialized incomplete (heuristic) algorithm
- General purpose incomplete algorithm

e.g.: simulated annealing, genetic algorithms, etc.

Divide and Concur belongs to the last category

Two constraint feasibility $x \in A \cap B$

Example:

A = permutations of "acgiknp"
B = 7-letter English words

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More structure

A, B are sets in a Euclidean configuration space Ω simple constraints: easy, efficient **projections** to A, B

$$P_A(x) = y \in A$$
 s.t. $||x-y||$ is minimized

Brief (incomplete) history of $x_{i+1} = x_i + P_B(2P_A(x_i) - x_i) - P_A(x_i)$

J. Douglas and H. H. Rachford, *On the numerical solution of heat conduction problems in two or three space variables*, Trans. Am. Math. Soc. 82 (1956), 421–439. splitting scheme for numerical PDE solutions

J.R. Fienup, Phase retrieval algorithms: a comparison, Applied Optics 21 (1982), 2758-2769. rediscovery, control theory motivation, phase retrieval

V. Elser, I. Rankenburg, and P. Thibault, Searching with iterated maps, PNAS 104, (2007), 418-423.

generalized form, applied to hard/frustrated problems: spin glass, SAT, protein folding, Latin squares, etc.

Projection to the packing (no overlaps) constraint



Projection to the packing (no overlaps) constraint



Dividing the Constraints



Dividing the Constraints



Dividing the Constraints



Projection to concurrence constraint



Projection to concurrence constraint



Divide and Concur scheme



No overlaps between designated replicas

All replicas of a particular figure concur

"divided" packing constraints "concurrence" constraint What can we do with projections?

• alternating projections:

$$x'_{i} = P_{A}(x_{i}); \quad x_{i+1} = P_{B}(x'_{i})$$

• Douglas-Rachford iteration (a/k/a difference map):

$$x_{i+1} = x_i + P_B(2P_A(x_i) - x_i) - P_A(x_i)$$

Finite packing problems



Gravel & Elser, Phys. Rev. E (2008)

Finite packing problems







Α

"divided" packing constraints (rigidity relaxed) В

"concurrence" + rigidity constraints



Α

"divided" packing constraints (rigidity relaxed) В

"concurrence" + rigidity constraints



Generalization to periodic packings



Generalization to periodic packings



Sphere packing and kissing in higher dimensions







"physical" tetrahedra



Kallus & Elser, preprint (2010)